



IBM Systems & Technology Group

z/VM Performance Introduction

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Overview

- Performance definition
- Guidelines
- Native CP commands
- Other performance tools
- I/O performance concepts
- Case study
- Final thoughts

Definition of Performance

Performance definitions:

- ☐ Response time
- ☐ Batch elapsed time
- ☐ Throughput
- ☐ Utilization
- ☐ Users supported
- ☐ Phone ringing
- ☐ Consistency
- ☐ All of the above

Performance Guidelines

- Processor
- Storage
- Paging
- Minidisk cache
- Server machines

Processor Guidelines

- Dedicated processors - mostly political
 - ▶ Absolute Share can be almost as effective
 - ▶ A virtual machine should have all dedicated or all shared processors
 - ▶ Gets wait state assist and 500ms minor time slice
- Share settings
 - ▶ Use absolute if you can judge percent of resources required
 - ▶ Use relative if difficult to judge and if lower share as system load increases is acceptable
 - ▶ Do not use LIMITHARD settings unnecessarily
- Do not define more virtual processors than are needed.
- If you use absolute LIMITHARD and need accuracy, use SET SRM LIMITHARD CONSUMPTION (requires VM64721)
- Small minor time slice keeps CP reactive.

Memory Guidelines

- Virtual:Real ratio should be $\leq 3:1$ or make sure you have robust paging system.
 - To avoid any performance impact for production workloads, you may need to keep ratio to 1:1
 - See also <http://www.vm.ibm.com/perf/tips/memory.html>
- Use SET RESERVE instead of LOCK to keep users pages in storage
- Define some processor storage as expanded storage to provide paging hierarchy For more background, see
 - <http://www.vm.ibm.com/perf/tips/storconf.html>
- Exploit shared memory where appropriate
- Size guests appropriately
 - Avoiding over provisioning
 - Do not put them in a high guest paging position

Paging Guidelines

- DASD paging allocations less than or equal to 50%.
 - QUERY ALLOC PAGE
- Watch blocks read per paging request (keep >10)
 - Monitor data (Performance Toolkit FCX103)
- Multiple volumes and multiple paths
- Do not mix Page extents with other extents on same volume.
- Paging volumes should all be of the same geometry and performance characteristics.
- Paging to FCP SCSI may offer higher paging bandwidth with higher processor requirements
- In a RAID environment, enable cache to mitigate write penalty.
- See also <http://www.vm.ibm.com/perf/tips/prgpage.html>

Minidisk Cache Guidelines

- Configure some real storage for MDC.
- In general, enable MDC for everything.
- Disable MDC for
 - ▶ Minidisks mapped to VM data spaces
 - ▶ write-mostly or read-once disks (logs, accounting)
 - ▶ Backup applications
- In large storage environments, may need to bias against MDC.
- Set maximum MDC limits
- Better performer than vdisks for read I/Os

Server Machine Guidelines

- Server Virtual Machine (SVM)
 - ▶ TCP/IP, RACFVM, etc.
- QUICKDSP ON to avoid eligible list
- Higher SHARE setting
- SET RESERVED to avoid paging
- NOMDCFS in directory option
- Ensure performance data includes these virtual machines

CP INDICATE Command

- **LOAD:** shows total system load.
 - ▶ (STORAGE value not very meaningful and was removed in z/VM 5.2.0)
- **USER EXP:** more useful than Indicate User
- **QUEUES EXP:** great for scheduler problems and quick state sampling
- **PAGING:** lists users in page wait.
- **IO:** lists users in I/O wait.
- **ACTIVE:** displays number of active users over given interval

CP INDICATE LOAD Example

INDICATE LOAD

AVGPROC-088% 03

XSTORE-000000/SEC MIGRATE-0000/SEC

MDC READS-000035/SEC WRITES-000001/SEC HIT RATIO-099%

PAGING-0023/SEC STEAL-000%

Q0-00007(00000)

DORMANT-00410

Q1-00000(00000)

E1-00000(00000)

Q2-00001(00000) EXPAN-002 E2-00000(00000)

Q3-00013(00000) EXPAN-002 E3-00000(00000)

PROC 0000-087%

PROC 0001-088%

PROC 0002-089%

LIMITED-00000

CP INDICATE QUEUE Example

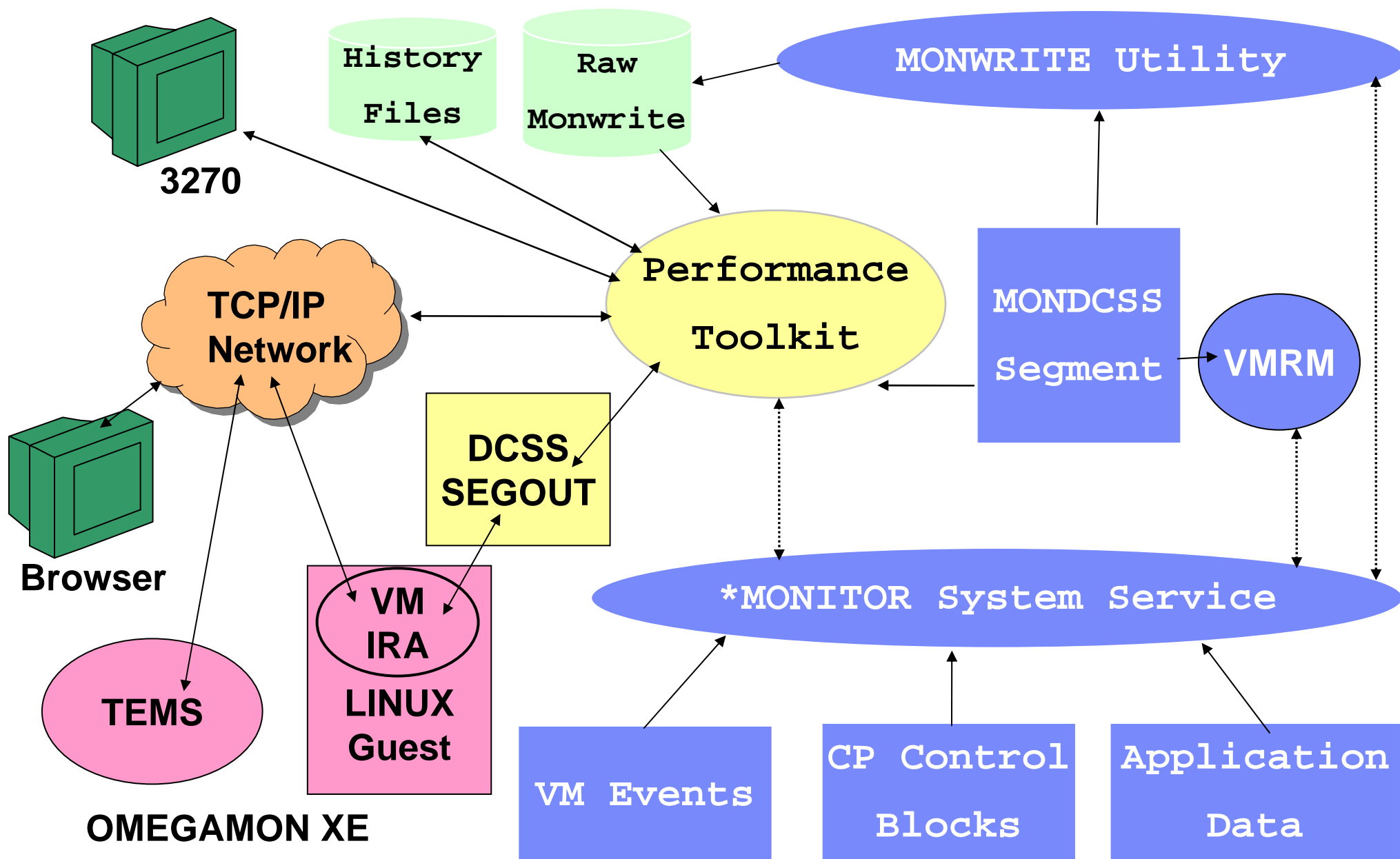
INDICATE QUEUE EXP

EDLLIB14	Q3	IO	00002473/00002654	..D.	-.0217	A00
KAZDAKC	Q3	IO	00003964/00003572	-.0190	A02
BITNER	Q1	R00	00001073/00001054	.I..	-.0163	A01
LCRAMER	Q3	IO	00003122/000028500259	A00
DSSERV	L0	R	00007290/000072893229	A00
RSCS	Q0	PS	00001638/00001616	.I..	99999	A00
SICIGANO	Q3	PS	00000662/00000662	.I..	99999	A00
VMLINUX1	Q3	PS	00018063/00018063	99999	A02
LNXREGR	Q3	PS	00073326/00073210	99999	A02
VMLINUX	Q3	PS	00031672/00031672	99999	A01
TCPIP	Q0	PS	00018863/00018397	.I..	99999	A02
EDLLNX2	Q3	PS	00032497/00032497	99999	A01
EDLLNX1	Q3	PS	00015939/00015939	99999	A02

Selected CP QUERY Commands

- Users: number and type of users on system
- SRM: scheduler/dispatcher settings
- SHARE: type and intensity of system share
- FRAMES: real storage allocation
- PATHS: physical paths to device and status
- ALLOC MAP: DASD allocation
- XSTORE: assignment of expanded storage
- MONITOR: current monitor settings
- MDC: MDC usage
- VDISK: virtual disk in storage usage
- SXSPAGES: System Execution Space (5.2.0)

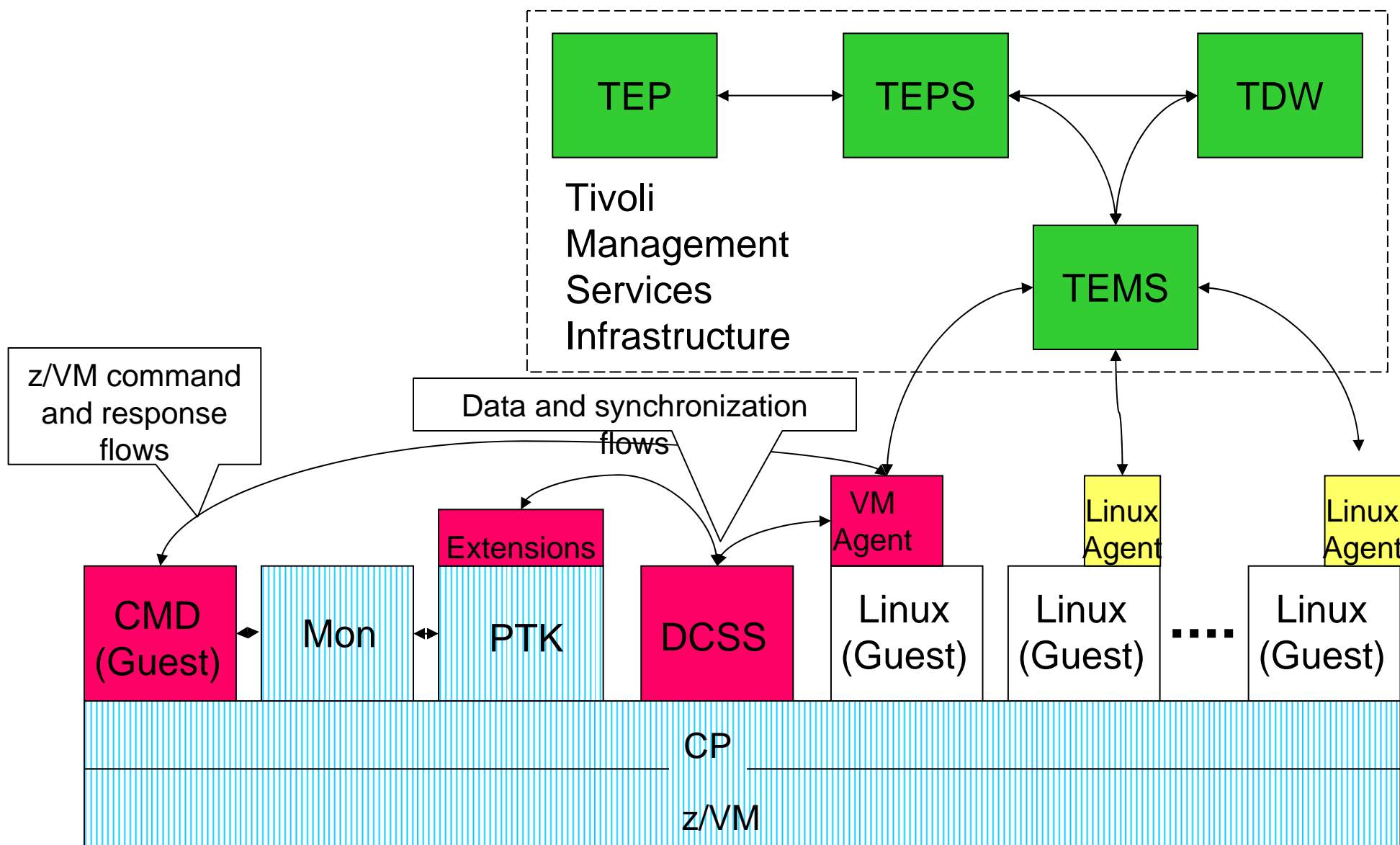
5,000 Foot View



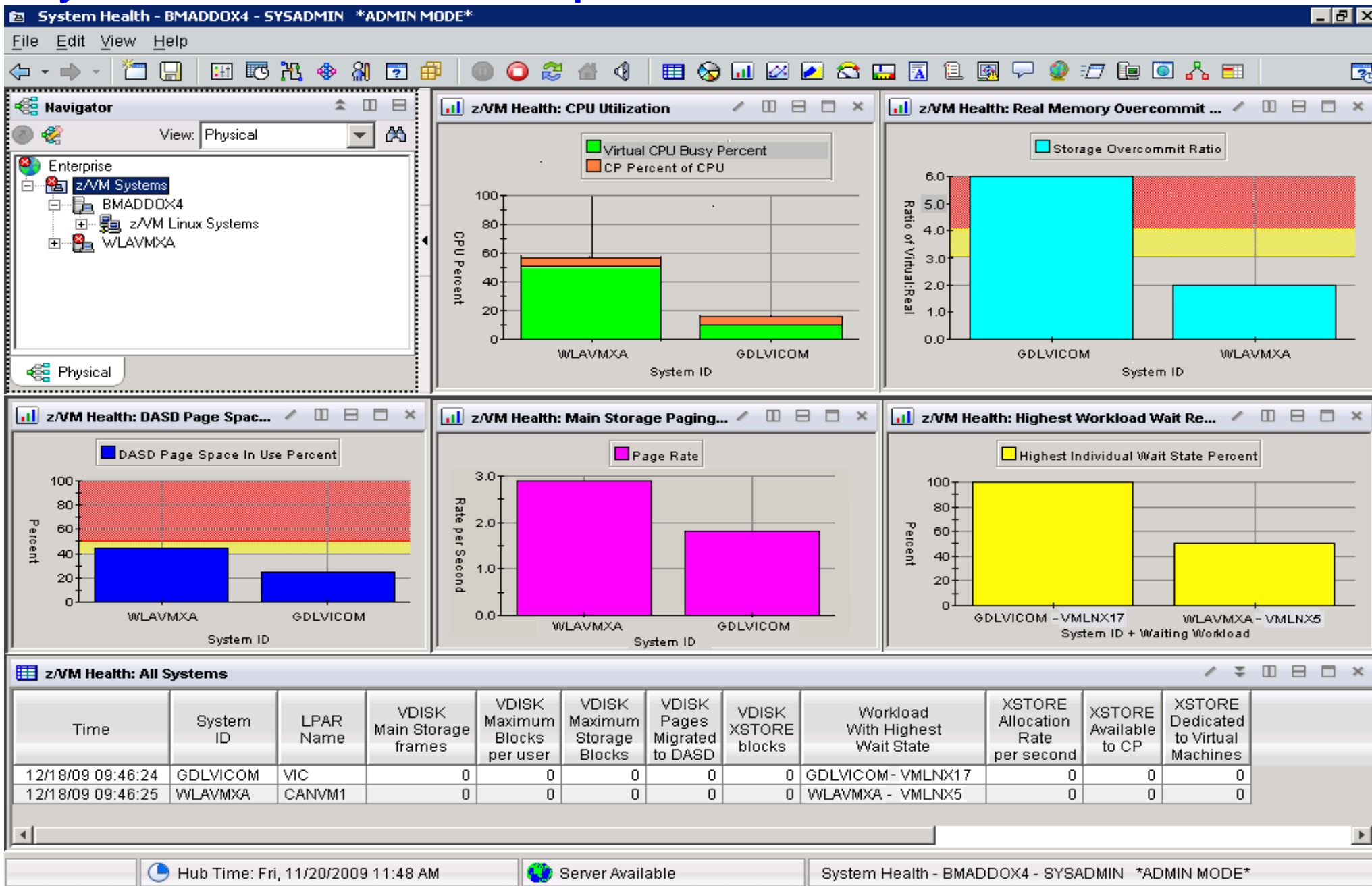
OMEGAMON XE on z/VM and Linux

- **Provides performance monitoring for z/VM and Linux guests**
- **Linux agents gather detailed performance data from Linux guests**
- **z/VM agent gathers performance data from z/VM and Linux**
 - Including z/VM view of guests
 - Uses IBM Performance Toolkit as its data source
- **Executes automated actions in response to defined events or situations**
- **Part of the Tivoli Management Services infrastructure and OMEGAMON family of products.**
 - Integrates performance management for your entire enterprise
- **Tivoli Data Warehouse (TDW) allows you to analyze trends and look at historical data**

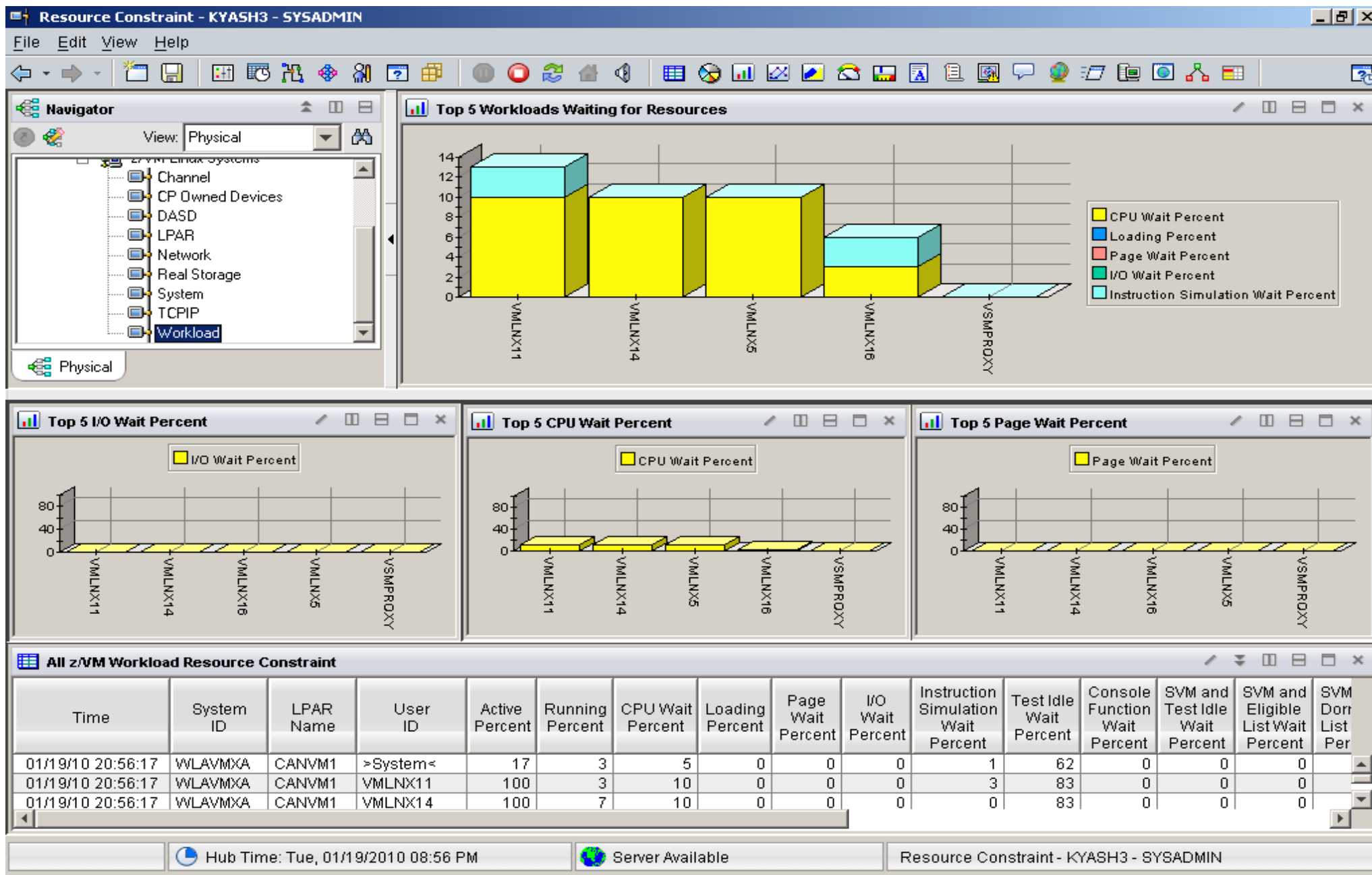
OMEGAMON XE Basic Architecture



System Health Workspace



V4.1.2 IF 1: Resource Constraint Analysis (Waits)



State Sampling

- Find the state of given user or device
 - ▶ Consolidation of samples gives useful info
- Low frequency:
 - ▶ INDICATE QUEUES
- High frequency:
 - ▶ Monitor: user, processor, and I/O domains
 - ▶ CP MONITOR SAMPLE RATE
- ▶ In Performance Toolkit
 - ▶ FCX108 DEVICE
 - ▶ FCX114 USTAT

I/O Response Time

$$\text{Resp Time} = \text{Service Time} + \text{Queue Time}$$

$$\text{Service Time} = \text{Pending} + \text{Connect} + \text{Disconnect}$$

- Queue Time: from hi-frequency sampling of queue in RDEV. Reported in monitor.
- Function Pending: time accumulated when a path to device cannot be obtained.
 - ▶ < 1 ms, unless contention at channels or control units.
- Connect: time device logically connected to channel path
 - ▶ proportional to amount of data per I/O

I/O Response Time (*continued*)

- Disconnect: time accumulated when device is logically disconnected from channel while subchannel system is active.
 - ▶ Cache miss
 - ▶ CU management
- Device Active: time accumulated between return of channel-end and device-end
 - ▶ Often reported as part of Disconnect Time

Definitions

- WSS = working set size
 - ▶ Comp-Sci Definition: Set of pages a workload needs to run effectively
 - ▶ VM Definition: Estimated working set size based primarily on resident page count
- Transaction
 - ▶ Comp-Sci Definition: End user interaction
 - ▶ VM Definition: transaction ends when scheduler detects end of processing

Other Sources

- Performance Manual - Part of z/VM Library
 - ▶ SC24-6208-00 z/VM 6.1.0
- <http://www.vm.ibm.com/perf/>
 - ▶ links to documents, tools, reference material
- <http://www.vm.ibm.com/perf/tips/>
 - ▶ common problems and solutions
 - ▶ guidelines
- <http://www.vm.ibm.com/devpages/bitner/>
 - ▶ presentations with speaker notes

A Case Study

The Grinch That Stole Performance

From Performance Toolkit USTAT FCX114 Report January 5:

<-SVM and->

%CPU	%LDG	%PGW	%IOW	%SIM	%TIW	%CFW	%TI	%EL	%DM	%IOA
0	0	0	19	2	10	0	3	0	51	8

From Performance Toolkit DEVICE FCX108 Report January 5:

<-Rate/s-> <----- Time (msec) -----> Req. <Pct>

Addr	I/O	Avoid	Pend	Disc	Conn	Serv	Resp	CUWt	Qued	Busy
1742	26.7	.0	1.3	18.4	4.7	24.5	69.0	.0	1.2	65.4

Went to check Toolkit CACHEXT FCX177 Report for control unit cache stats, but it didn't exist!

It is a good thing I keep historical data -- let's go back and see what's going on...

When Did We Last See Cache?

From Performance Toolkit DEVICE FCX108 Report:

	<-Rate/s->	<----- Time (msec) ----->						Req.	<Pct>	
Addr	I/O	Avoid	Pend	Disc	Conn	Serv	Resp	CUWt	Qued	Busy
Dec8	41.0	.0	0.3	0.2	2.0	2.6	2.9	.0	.0	10.5
Jan5	26.7	.0	1.3	18.4	4.7	24.5	69.0	.0	1.2	65.4

From Performance Toolkit CACHEXT FCX177 Dec. 8th Report:

<----- Rate/s ----->					<-----Percent----->					
Total	Total	Read	Read	Write	<----- Hits ----->					
Cache	SCMBK	N-Seq	Seq	FW	Read	Tot	RdHt	Wrt	DFW	CFW
53.0	41.0	52.3	0	0.6	99	99	99	96	96	..

Down for the 3-Count

```
q dasd details 1742
```

```
1742 CUTYPE = 3990-EC, DEVTYPE = 3390-06,  
VOLSER=USE001
```

```
    CACHE DETAILS:  CACHE NVS  CFW  DFW  PINNED  CONCOPY  
      -SUBSYSTEM    F      Y    Y    -    Y          N  
      -DEVICE       Y      -    -    Y    N          N  
DEVICE DETAILS:  CCA = 02, DDC = 02  
DUPLEX DETAILS:  SIMPLEX
```

Pinned data! Yikes! I had never seen that before!

Performance Toolkit Device Details

FCX110 CPU 2003 GDLVM7 Interval INITIAL. - 13:08:47 Remote Data

Detailed Analysis for Device 1742 (SYSTEM)

Device type :	3390-2	Function pend.:	.8ms	Device busy :	27%
VOLSER :	USE001	Disconnected :	20.3ms	I/O contention:	0%
Nr. of LINKs:	404	Connected :	5.4ms	Reserved :	0%
Last SEEK :	1726	Service time :	26.5ms	SENSE SSCH :	...
SSCH rate/s :	10.5	Response time :	26.5ms	Recovery SSCH :	...
Avoided/s :	CU queue time :	.0ms	Throttle del/s:	...

Status: SHARABLE

Path(s) to device 1742: 0A 2A 4A

Channel path status : ON ON ON

Device	Overall CU-Cache Performance								Split
DIR ADDR VOLSER	IO/S	%READ	%RDHIT	%WRHIT	ICL/S	BYP/S	IO/S	%READ	%RDHIT
08 1742 USE001	.0	0	0	0	.0	.0	'NORMAL'	I/O only	

Performance Toolkit Device Details

MDISK	Extent	Userid	Addr	Status	LINK	MDIO/s

101	- 200	EDLSFS	0310	WR	1	.0
201	- 500	EDLSFS	0300	WR	1	.0
501	- 600	EDLSFS	0420	WR	1	.0
601	- 1200	EDLSFS	0486	WR	1	.0
1206	- 1210	RAID	0199	owner		
		BRIANKT	0199	RR	5	.0
1226	- 1525	DATABASE	0465	owner		
		K007641	03A0	RR	3	.0
1526	- 1625	DATABASE	0269	owner		
		BASILEMM	0124	RR	25	.0
1626	- 1725	DATABASE	0475	owner		
		SUSANF7	0475	RR	1	.0
1726	- 2225	DATABASE	0233	owner	366	10.5

Solution

- Use **Q PINNED** CP command to check for what data is pinned.
- Discussion with Storage Management team.
- Moved data off string until corrected.

Pinned data is very rare, but when it happens it is serious.

Some Final Thoughts

- Collect data for a base line of good performance.
- Implement change management process.
- Make as few changes as possible at a time.
- Performance is often only as good as the weakest component.
- Relieving one bottleneck will reveal another. As attributes of one resource change, expect at least one other to change as well.
- Latent demand is real.